

# Probabilistic Context-free Grammars

Computational Linguistics

Alexander Koller

22 November 2019

# The CKY Recognizer

$S \rightarrow NP VP$

$V \rightarrow ate$

$Det \rightarrow a$

$NP \rightarrow Det N$

$NP \rightarrow John$

$N \rightarrow sandwich$

$VP \rightarrow V NP$

Chart

$S \Rightarrow^* w$

	$i = 1$	2	3	4
5	S	VP	NP	N
4			Det	... a sandwich
3		V	... ate	a
$k = 2$	NP	... John	ate	
	John			

Cell at column  $i$ , row  $k$ :  
 $\{ A \mid A \Rightarrow^* w_i \dots w_{k-1} \}$

# CKY recognizer: pseudocode

Data structure:  $\text{Ch}(i,k)$  eventually contains  $\{A \mid A \Rightarrow^* w_i \dots w_{k-1}\}$   
(initially all empty).

for each  $i$  from 1 to  $n$ :

  for each production rule  $A \rightarrow w_i$ :

    add  $A$  to  $\text{Ch}(i, i+1)$

for each *width*  $b$  from 2 to  $n$ :

  for each *start position*  $i$  from 1 to  $n-b+1$ :

    for each *left width*  $k$  from 1 to  $b-1$ :

      for each  $B \in \text{Ch}(i, i+k)$  and  $C \in \text{Ch}(i+k, i+b)$ :

        for each production rule  $A \rightarrow B C$ :

          add  $A$  to  $\text{Ch}(i, i+b)$

claim that  $w \in L(G)$  iff  $S \in \text{Ch}(1, n+1)$

# Recognizer to Parser

$S \rightarrow NP VP$

$V \rightarrow \text{ate}$

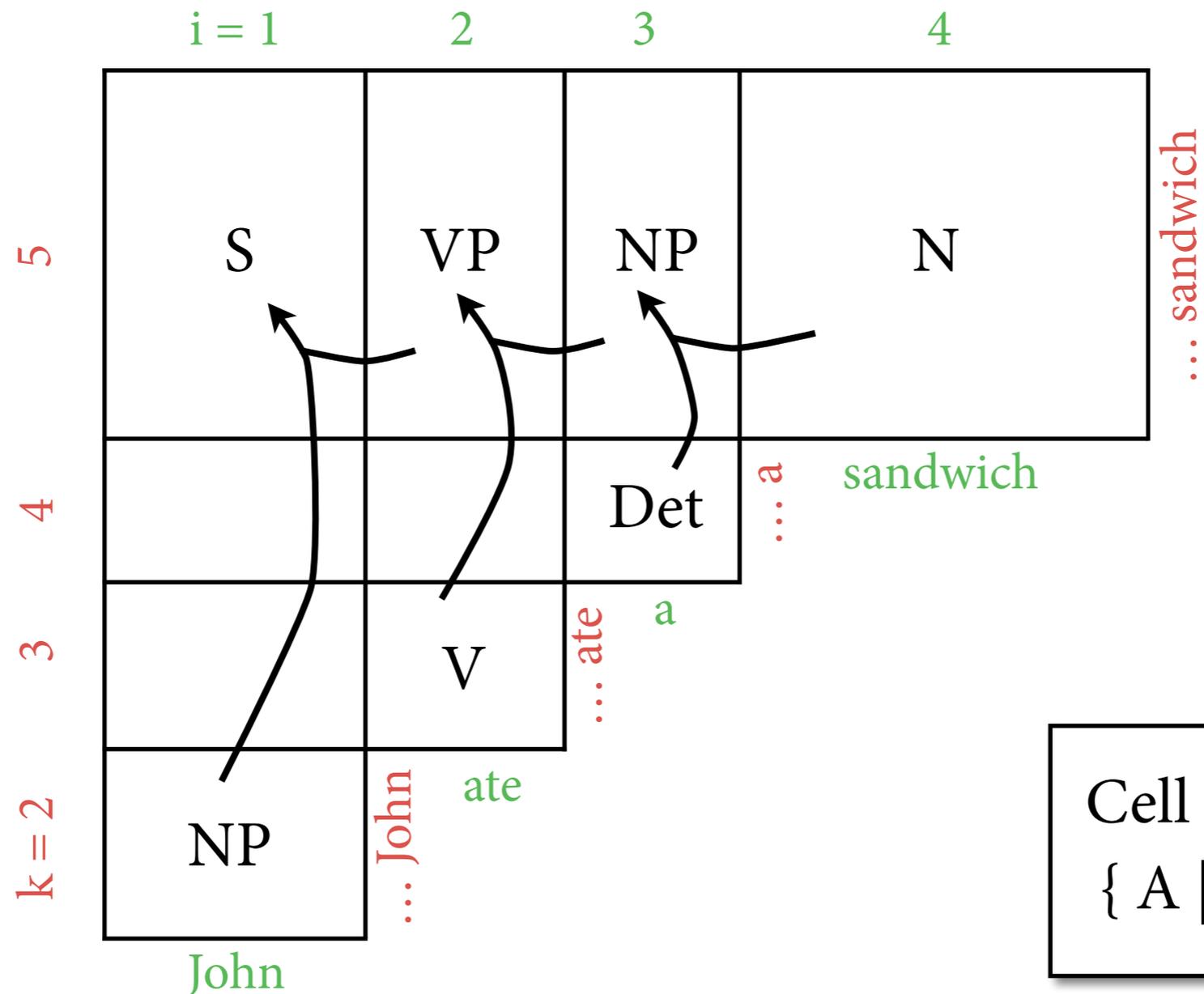
$\text{Det} \rightarrow \text{a}$

$NP \rightarrow \text{Det } N$

$NP \rightarrow \text{John}$

$N \rightarrow \text{sandwich}$

$VP \rightarrow V NP$



Cell at column  $i$ , row  $k$ :  
 $\{ A \mid A \Rightarrow^* w_i \dots w_{k-1} \}$

# Recognizer to Parser

- Parser: need to construct parse trees from chart.
- Do this by memorizing how each  $A \in \text{Ch}(i,k)$  can be constructed from smaller parts.
  - ▶ built from  $B \in \text{Ch}(i,j)$  and  $C \in \text{Ch}(j,k)$  using  $A \rightarrow B C$ : store  $(B,C,j)$  in *backpointer* for  $A$  in  $\text{Ch}(i,k)$ .
  - ▶ analogous to backpointers in HMMs
- Once chart has been filled, enumerate trees recursively by following backpointers, starting at  $S \in \text{Ch}(1,n+1)$ .

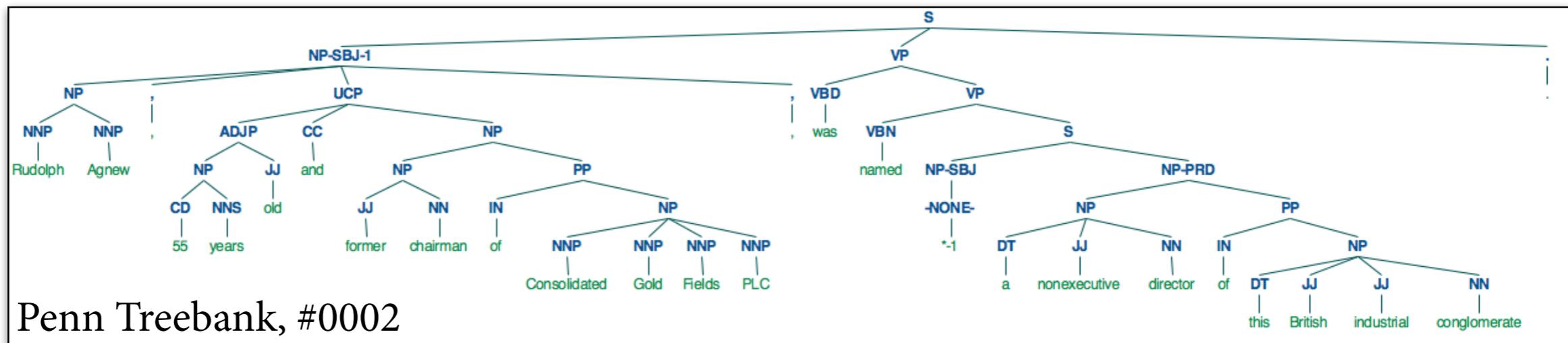
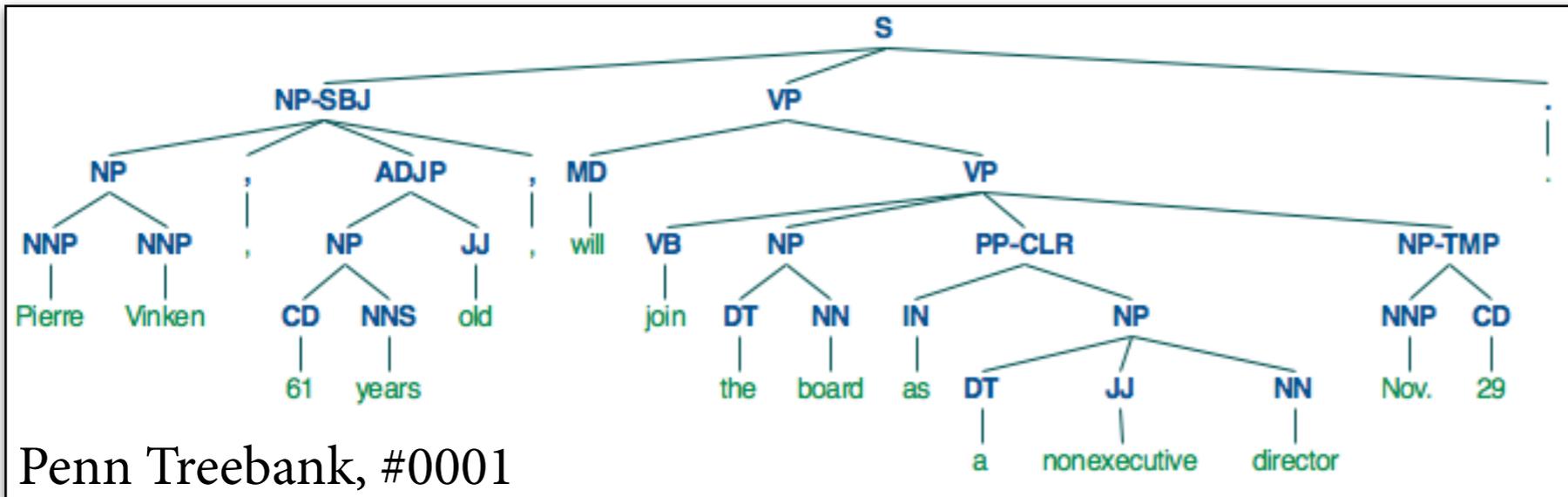
# Let's play a game

- Given a nonterminal symbol, expand it.
- You can take one of two moves:
  - ▶ expand nonterminal into a sequence of other nonterminals
  - ▶ use nonterminals S, NP, VP, PP, ... or POS tags
  - ▶ expand nonterminal into a word

# Penn Treebank POS tags

Tag	Description	Example	Tag	Description	Example
CC	Coordin. Conjunction	<i>and, but, or</i>	SYM	Symbol	<i>+, %, &amp;</i>
CD	Cardinal number	<i>one, two, three</i>	TO	“to”	<i>to</i>
DT	Determiner	<i>a, the</i>	UH	Interjection	<i>ah, oops</i>
EX	Existential ‘there’	<i>there</i>	VB	Verb, base form	<i>eat</i>
FW	Foreign word	<i>mea culpa</i>	VBD	Verb, past tense	<i>ate</i>
IN	Preposition/sub-conj	<i>of, in, by</i>	VBG	Verb, gerund	<i>eating</i>
JJ	Adjective	<i>yellow</i>	VBN	Verb, past participle	<i>eaten</i>
JJR	Adj., comparative	<i>bigger</i>	VBP	Verb, non-3sg pres	<i>eat</i>
JJS	Adj., superlative	<i>wildest</i>	VBZ	Verb, 3sg pres	<i>eats</i>
LS	List item marker	<i>1, 2, One</i>	WDT	Wh-determiner	<i>which, that</i>
MD	Modal	<i>can, should</i>	WP	Wh-pronoun	<i>what, who</i>
NN	Noun, sing. or mass	<i>llama</i>	WP\$	Possessive wh-	<i>whose</i>
NNS	Noun, plural	<i>llamas</i>	WRB	Wh-adverb	<i>how, where</i>
NNP	Proper noun, singular	<i>IBM</i>	\$	Dollar sign	<i>\$</i>
NNPS	Proper noun, plural	<i>Carolinas</i>	#	Pound sign	<i>#</i>
PDT	Predeterminer	<i>all, both</i>	“	Left quote	<i>( ‘ or “)</i>
POS	Possessive ending	<i>'s</i>	”	Right quote	<i>( ’ or ”)</i>
PP	Personal pronoun	<i>I, you, he</i>	(	Left parenthesis	<i>( [ , { , &lt;)</i>
PP\$	Possessive pronoun	<i>your, one's</i>	)	Right parenthesis	<i>( ] , } , &gt;)</i>
RB	Adverb	<i>quickly, never</i>	,	Comma	<i>,</i>
RBR	Adverb, comparative	<i>faster</i>	.	Sentence-final punc	<i>( . ! ?)</i>
RBS	Adverb, superlative	<i>fastest</i>	:	Mid-sentence punc	<i>( : ; ... - -)</i>
RP	Particle	<i>up, off</i>			

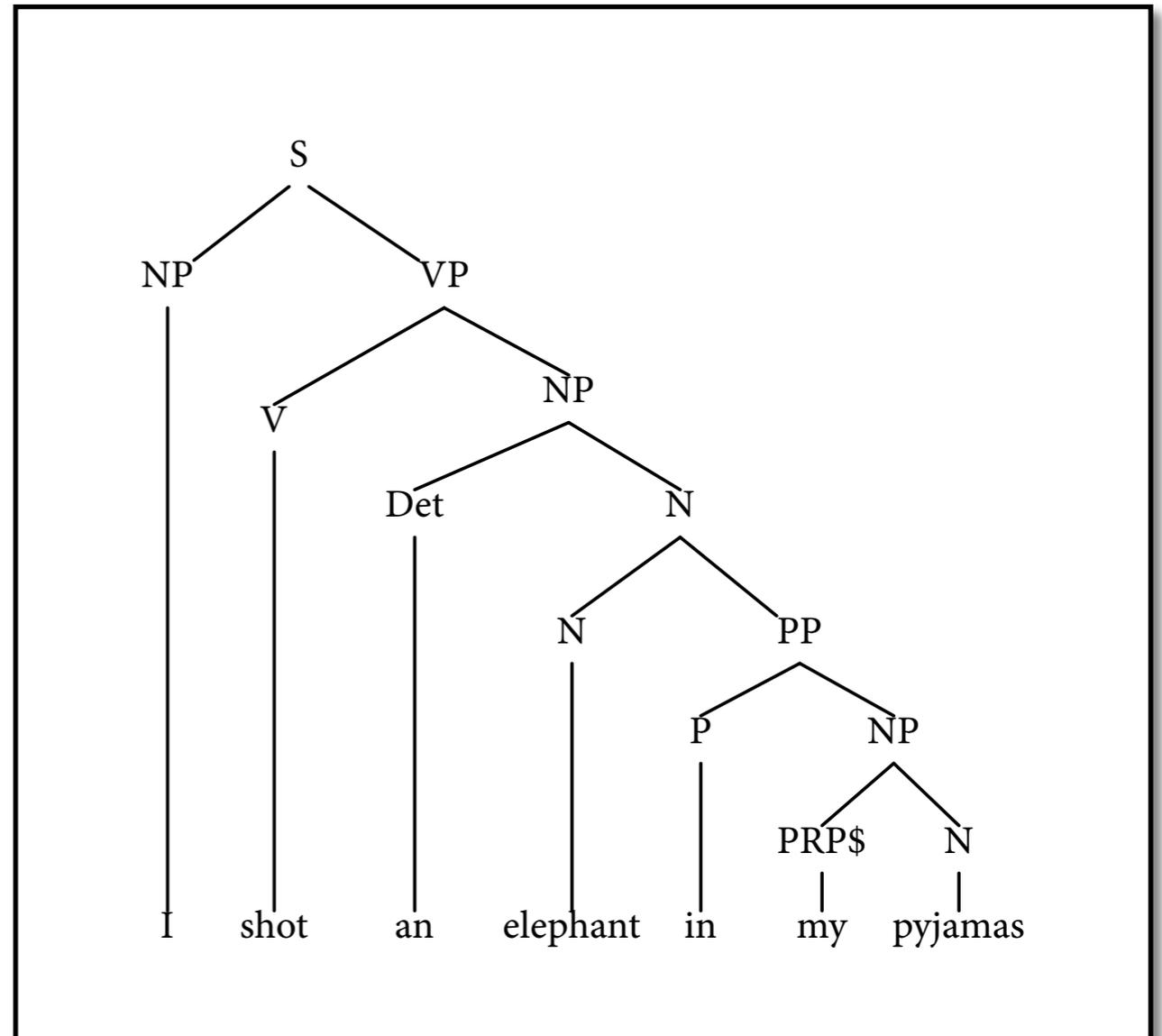
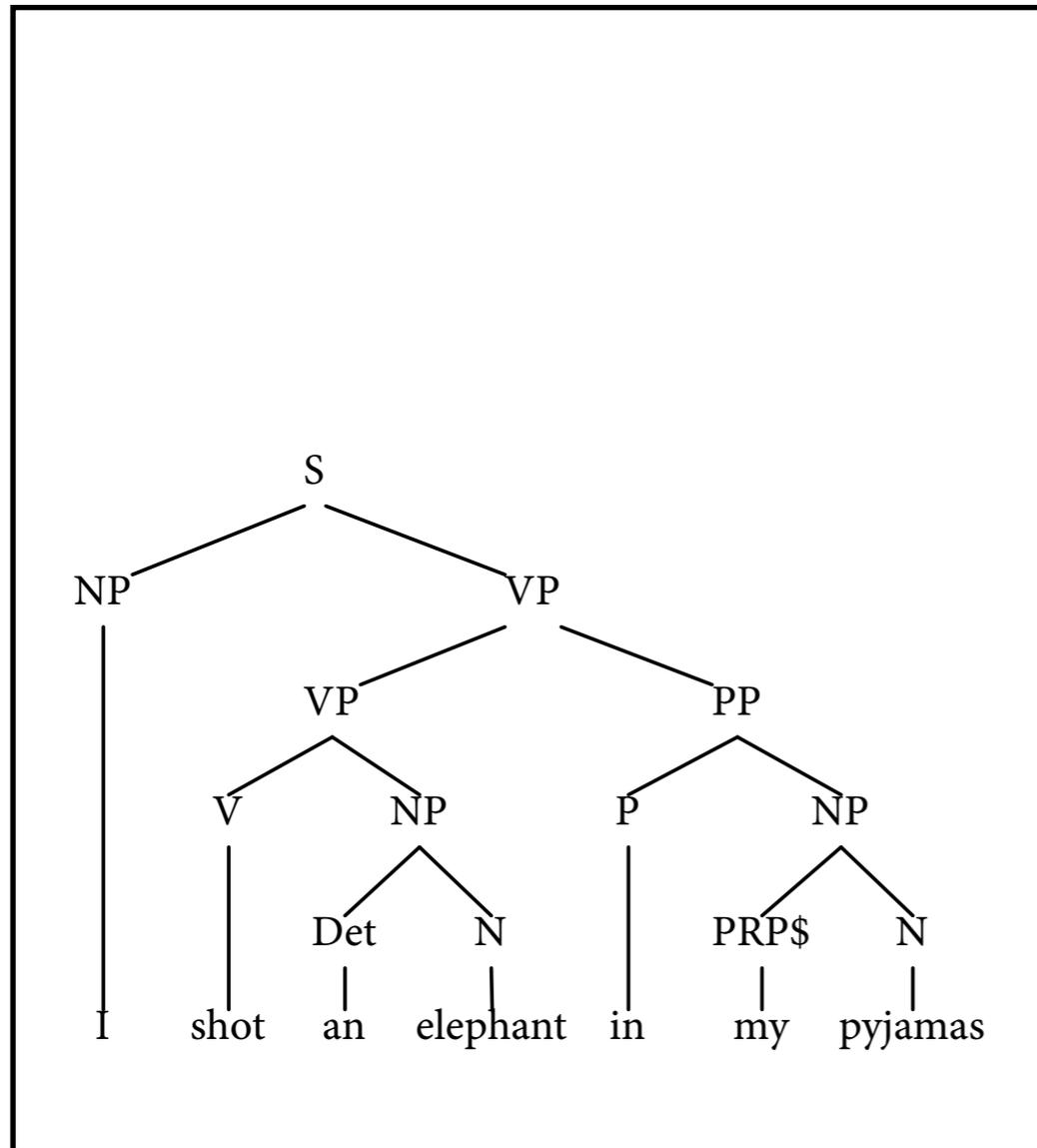
# Some real trees



`nltk.corpus.treebank.parsed_sents("wsj_0001.mrg")[0].draw()`

# Ambiguity

Need to *disambiguate*: find “correct” parse tree for ambiguous sentence.



How do we identify the “correct” tree?

How do we compute it efficiently? (Remember: exponential number of readings.)

# Probabilistic CFGs

- A *probabilistic context-free grammar (PCFG)* is a context-free grammar in which
  - ▶ each production rule  $A \rightarrow w$  has a probability  $P(A \rightarrow w \mid A)$ : when we expand  $A$ , how likely is it that we choose  $A \rightarrow w$ ?
  - ▶ for each nonterminal  $A$ , probabilities must sum to one:

$$\sum_w P(A \rightarrow w \mid A) = 1$$

- ▶ we will write  $P(A \rightarrow w)$  instead of  $P(A \rightarrow w \mid A)$  for short

# An example

$S \rightarrow NP VP$	[1.0]	$VP \rightarrow V NP$	[0.5]
$NP \rightarrow Det N$	[0.8]	$VP \rightarrow VP PP$	[0.5]
$NP \rightarrow i$	[0.2]	$V \rightarrow shot$	[1.0]
$N \rightarrow N PP$	[0.4]	$PP \rightarrow P NP$	[1.0]
$N \rightarrow elephant$	[0.3]	$P \rightarrow in$	[1.0]
$N \rightarrow pyjamas$	[0.3]	$Det \rightarrow an$	[0.5]
		$Det \rightarrow my$	[0.5]

(let's pretend for simplicity that Det = PRP\$)

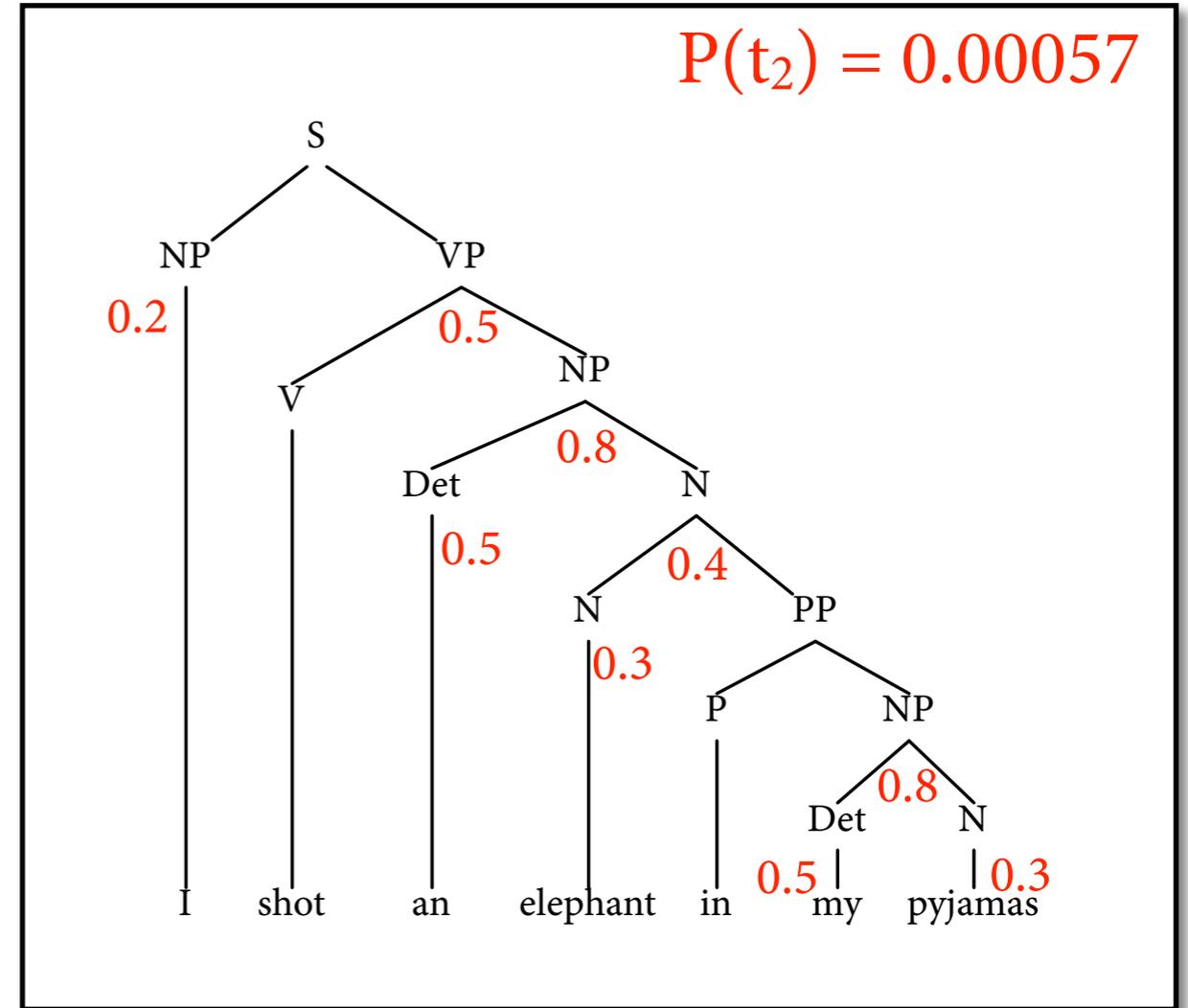
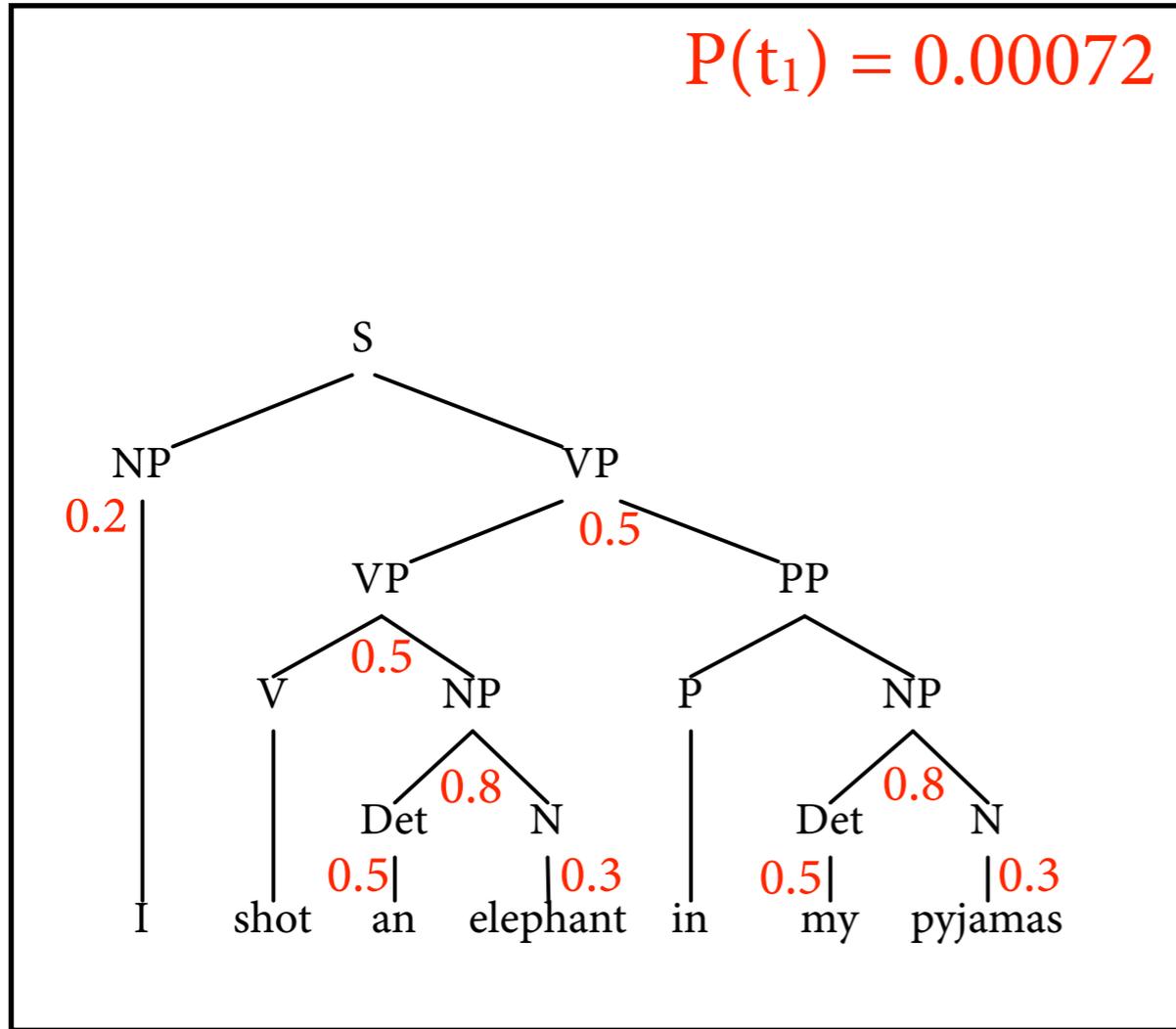
# Generative process

- PCFG generates random derivations of CFG.
  - ▶ each event (expand nonterminal by production rule) statistically independent of all the others

$S \xRightarrow{1.0} NP VP \xRightarrow{0.2} i VP \xRightarrow{0.5} i VP PP$   
 $\Rightarrow^* i \text{ shot an elephant in my pyjamas}$  0.00072

$S \xRightarrow{1.0} NP VP \xRightarrow{0.2} i VP \xRightarrow{0.4} i V Det N$   
 $\Rightarrow i V Det N PP \Rightarrow^* i \text{ shot ... pyjamas}$  0.00057  
0.4

# Parse trees



“correct” = more probable parse tree

# Language modeling

- As with other generative models (e.g. HMMs), can define probability  $P(w)$  of string by marginalizing over its possible parses:

$$P(w) = \sum_{t \in \text{parses}(w)} P(t)$$

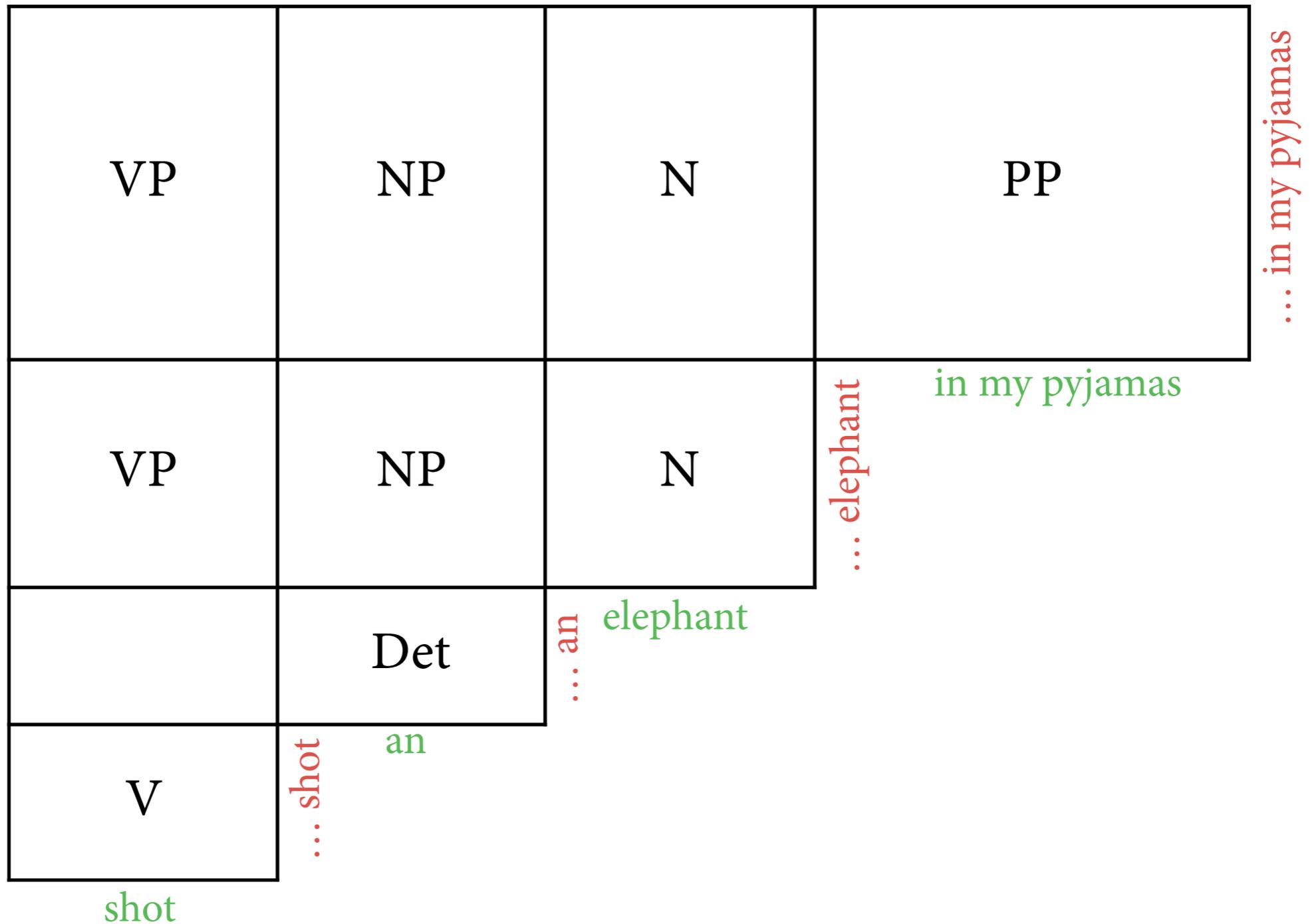
- Can compute this efficiently with *inside probabilities*, see next time.

# Disambiguation

- Assumption: “correct” parse tree = the parse tree that had highest prob of being generated by random process, i.e.  $\operatorname{argmax}_{t \in \text{pases}(w)} P(t)$
- We use a variant of the Viterbi algorithm to compute it.
- Here, Viterbi based on CKY; can do it with other parsing algorithms too.

# The intuition

Ordinary CKY parse chart:  $Ch(i,k) = \{A \mid A \Rightarrow^* w_i \dots w_{k-1}\}$



# The intuition

Viterbi CKY parse chart:  $\text{Ch}(i, k) = \{(A, p) \mid p = \max_{d: A \Rightarrow^* w_i \dots w_{k-1}} P(d)\}$

VP: 0.0036	NP: 0.006	N: 0.014	PP: 0.12	... in my pyjamas
VP: 0.06	NP: 0.12	N: 0.3	... elephant	in my pyjamas
	Det: 0.5	... an	elephant	
V: 1.0	... shot	an		
shot				

# Viterbi CKY

- Define for each span  $(i,k)$  and each nonterminal  $A$  the probability

$$V(A, i, k) = \max_{A \xRightarrow{d} w_i \dots w_{k-1}} P(d)$$

- Compute  $V$  iteratively “bottom up”, i.e. starting from small spans and working our way up to longer spans.

$$V(A, i, i + 1) = P(A \rightarrow w_i)$$

$$V(A, i, k) = \max_{\substack{A \rightarrow B C \\ i < j < k}} P(A \rightarrow B C) \cdot V(B, i, j) \cdot V(C, j, k)$$

# Viterbi CKY - pseudocode

set all  $V[A, i, j]$  to  $\emptyset$

for all  $i$  from 1 to  $n$ :

  for all  $A$  with rule  $A \rightarrow w_i$ :

    add  $A$  to  $Ch(i, i+1)$

$V[A, i, i+1] = P(A \rightarrow w_i)$

for all  $b$  from 2 to  $n$ :

  for all  $i$  from 1 to  $n-b+1$ :

    for all  $k$  from 1 to  $b-1$ :

      for all  $B$  in  $Ch(i, i+k)$  and  $C$  in  $Ch(i+k, i+b)$ :

        for all production rules  $A \rightarrow B C$ :

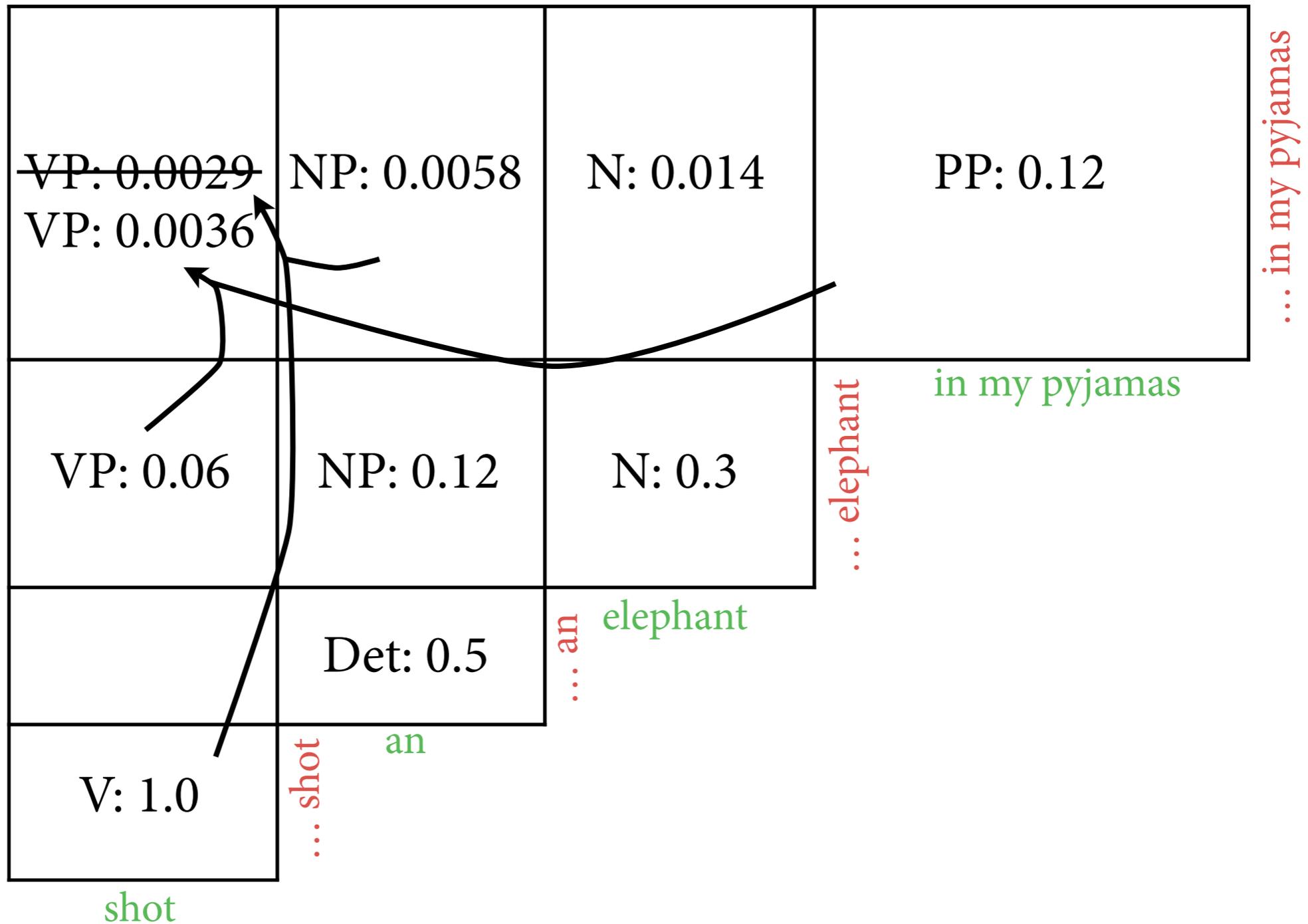
          add  $A$  to  $Ch(i, i+b)$

          if  $P(A \rightarrow B C) * V[B, i, i+k] * V[C, i+k, i+b] > V[A, i, i+b]$ :

$V[A, i, i+b] = P(A \rightarrow B C) * V[B, i, i+k] * V[C, i+k, i+b]$

# Viterbi-CKY in action

Viterbi CKY parse chart:  $Ch(i, k) = \{(A, p) \mid p = \max_{d:A \Rightarrow^* w_i \dots w_{k-1}} P(d)\}$



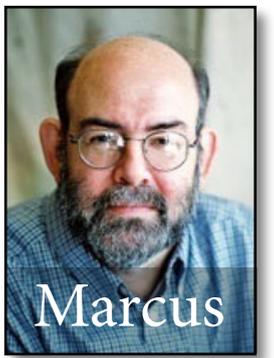
# Remarks

- Viterbi CKY has exactly the same nested loops as the ordinary CKY parser.
  - ▶ computing  $V$  in addition to  $Ch$  only changes constant factor
  - ▶ thus asymptotic runtime remains  $O(n^3)$
- Compute optimal parse by storing backpointers.
  - ▶ same backpointers as in ordinary CKY
  - ▶ sufficient to store the *best* backpointer for each  $(A,i,k)$  if we only care about best parse (and not all parses), i.e. actually uses less memory than ordinary CKY

# Obtaining the PCFG

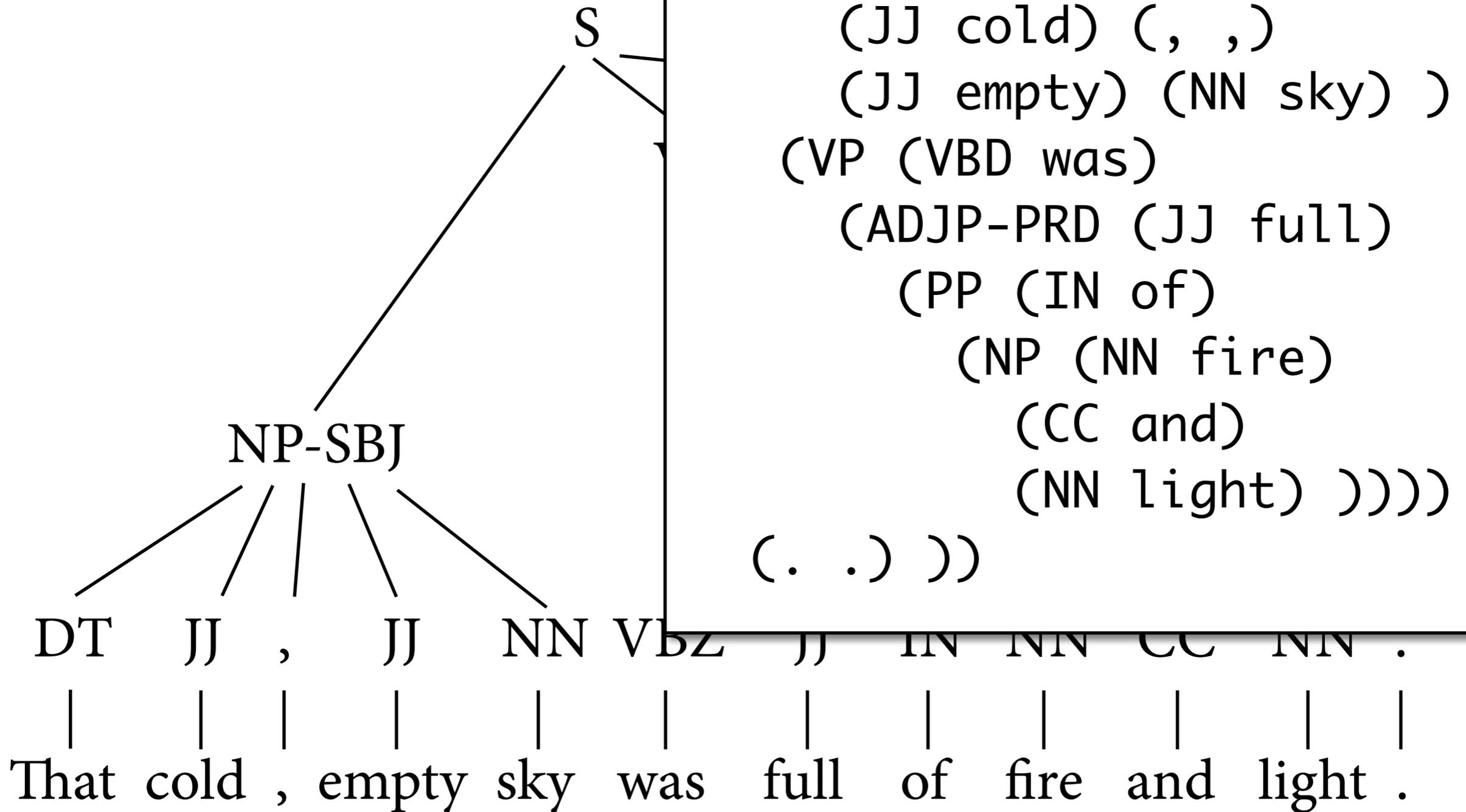
- How to obtain the CFG?
  - ▶ write by hand
  - ▶ derive from *treebank*
  - ▶ *grammar induction* from raw text
- How to obtain the rule probabilities once we have the CFG?
  - ▶ maximum likelihood estimation from treebank
  - ▶ EM training from raw text (inside-outside algorithm)

# The Penn Treebank



- Large (in the mid-90s) quantity of text, annotated with POS tags and syntactic structures.
- Consists of several sub-corpora:
  - ▶ Wall Street Journal: 1 year of news text, 1 million words
  - ▶ Brown corpus: balanced corpus, 1 million words
  - ▶ ATIS: dialogues on flight bookings, 5000 words
  - ▶ Switchboard: spoken dialogue, 3 million words
- WSJ PTB is standard corpus for training and evaluating PCFG parsers.

# Annotation format



```
((S
  (NP-SBJ (DT That)
    (JJ cold) (, ,)
    (JJ empty) (NN sky) )
  (VP (VBD was)
    (ADJP-PRD (JJ full)
      (PP (IN of)
        (NP (NN fire)
          (CC and)
          (NN light) ))))
    (. .) ))
```

# Reading off grammar

- Can directly read off “grammar in annotators’ heads” from trees in treebank.

- Yields very large CFG, e.g. 4500 rules for VP:

VP → VBD PP

VP → VBD PP PP

VP → VBD PP PP

VP → VBD PP PP

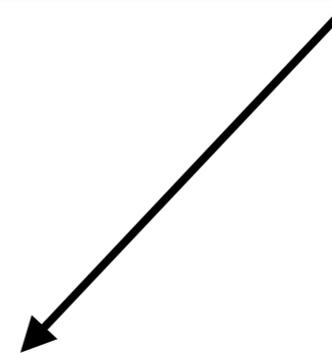
VP → VBD ADVP PP

VP → VBD PP ADVP

...

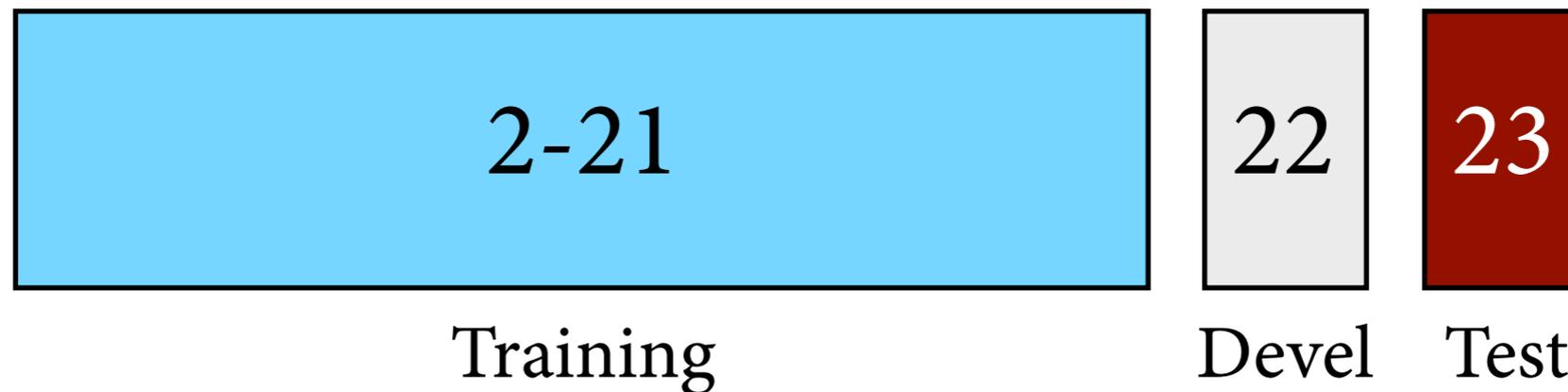
VP → VBD PP PP PP PP PP ADVP PP

“This mostly happens because we go from football in the fall to lifting in the winter to football again in the spring.”



# Evaluation

- Step 1: Decide on training and test corpus.  
For WSJ corpus, there is a conventional split by sections:

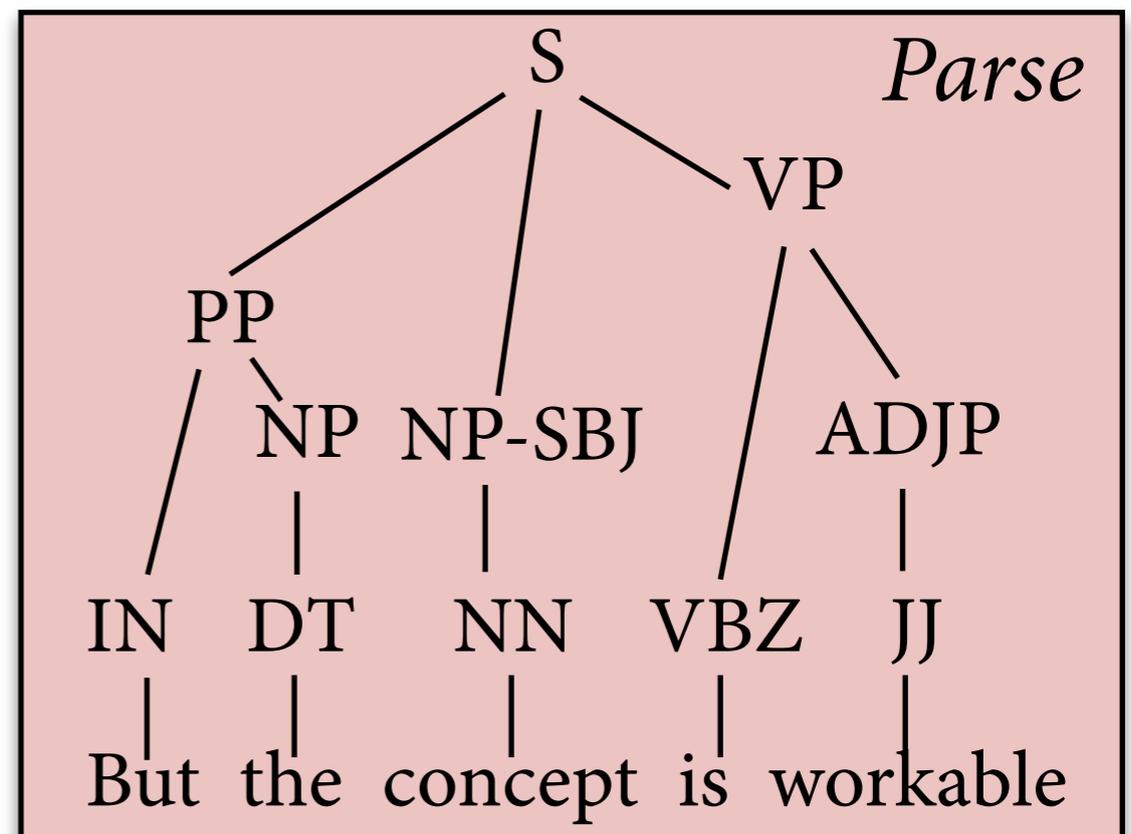
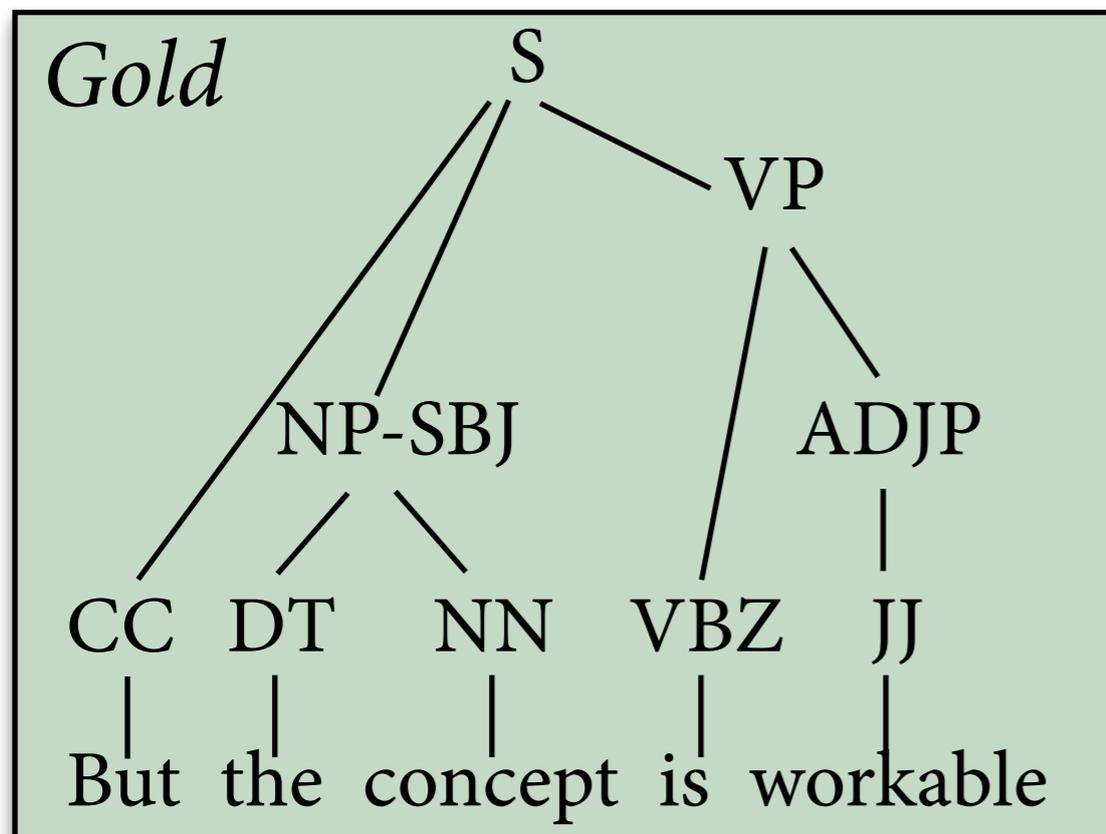


# Evaluation

- Step 2: How should we measure the accuracy of the parser?
- Straightforward idea: Measure “exact match”, i.e. proportion of gold standard trees that parser got right.
- This is too strict:
  - ▶ parser makes many decisions in parsing a sentence
  - ▶ a single incorrect parsing decision makes tree “wrong”
  - ▶ want more fine-grained measure

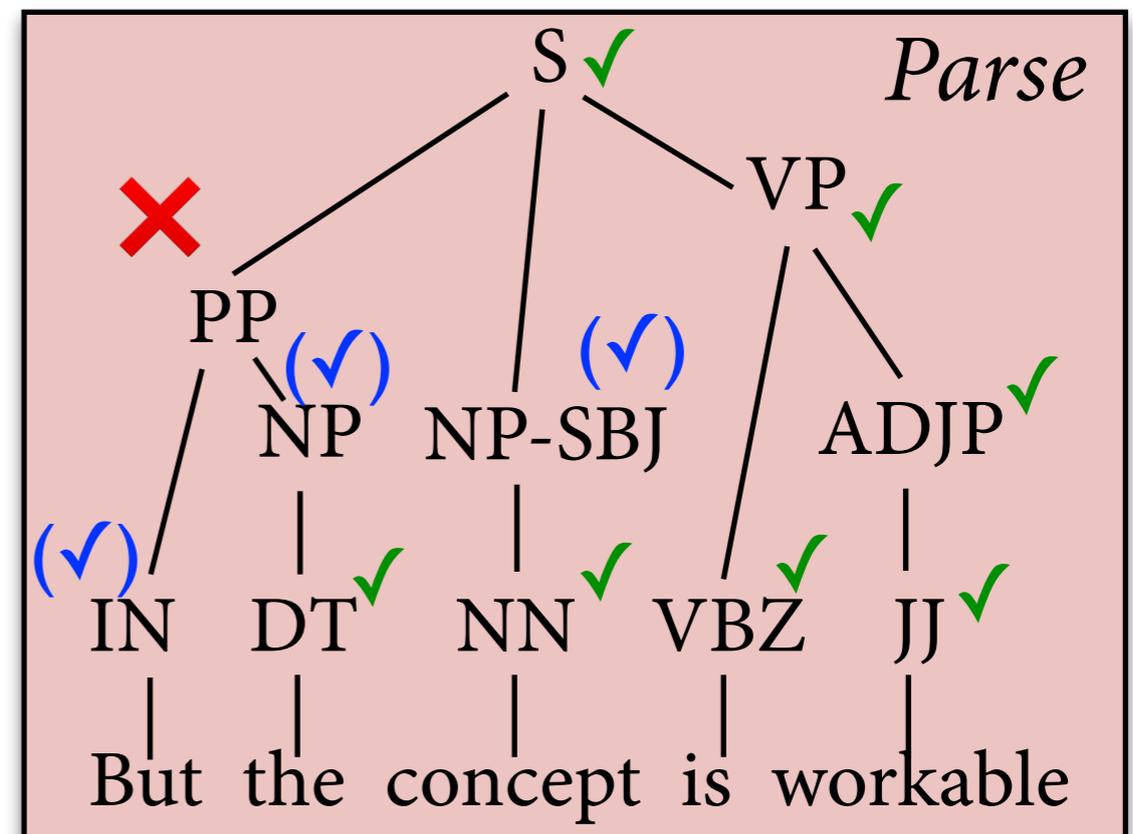
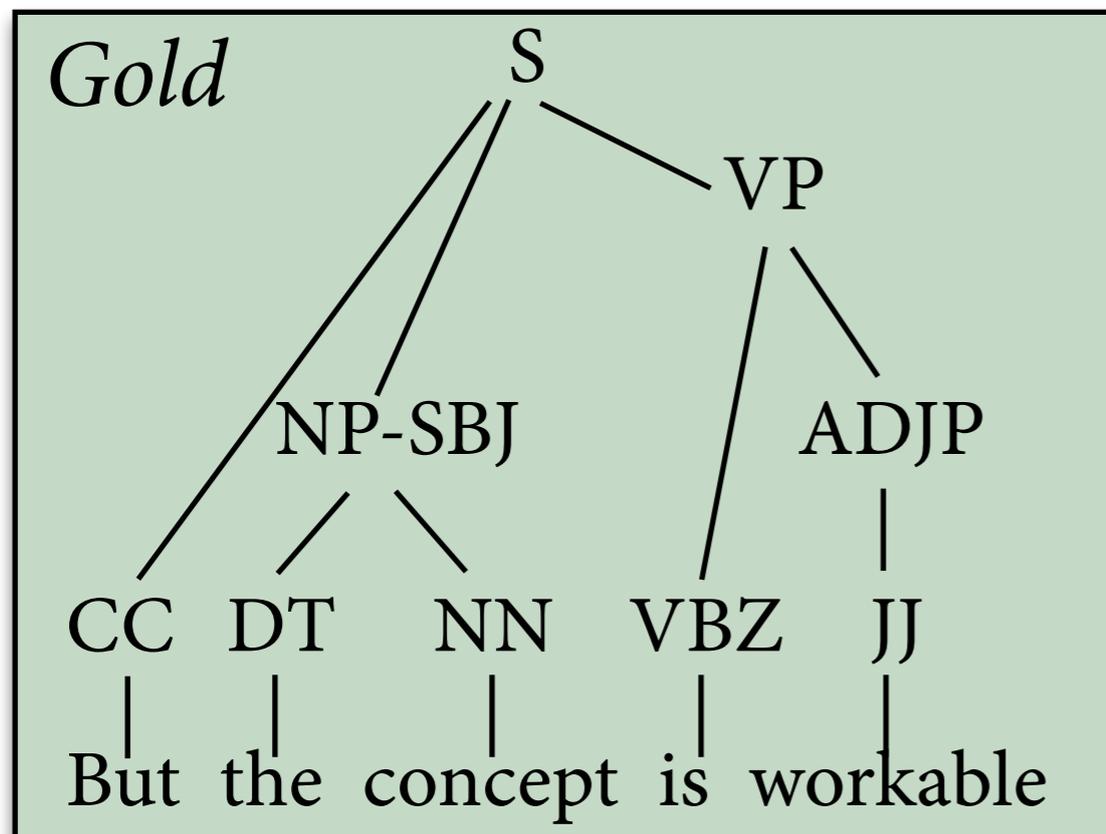
# Comparing parse trees

- Idea 2 (PARSEVAL): Compare *structure* of parse tree and gold standard tree.
  - ▶ Labeled: Which *constituents* (span + syntactic category) of one tree also occur in the other?
  - ▶ Unlabeled: How do the trees bracket the *substrings* of the sentence (ignoring syntactic categories)?



# Precision

What proportion of constituents in *parse tree* is also present in *gold tree*?

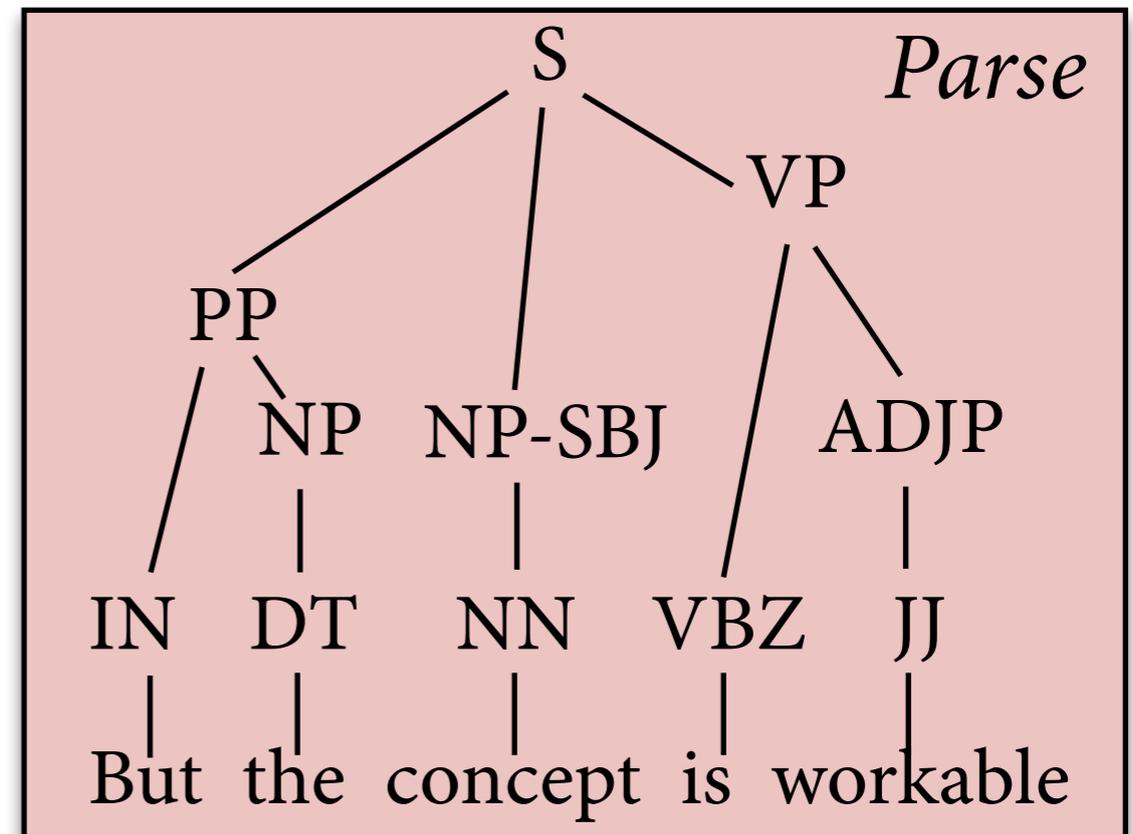
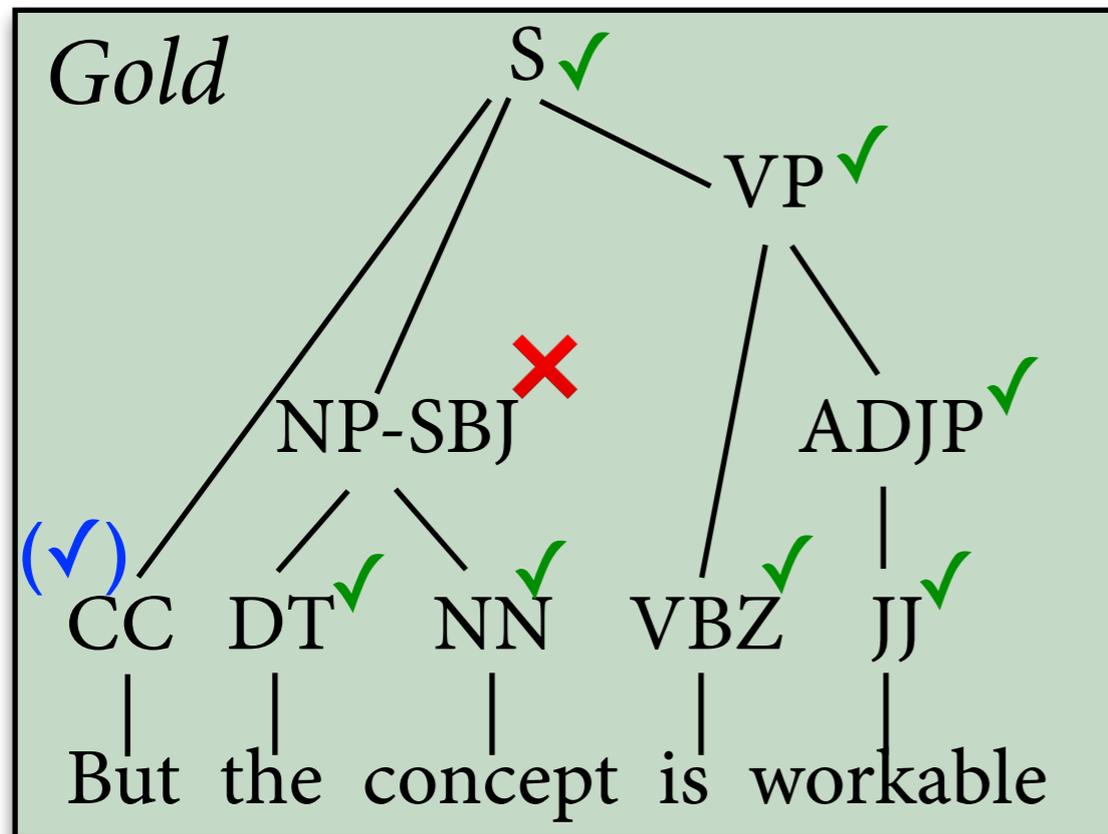


Labeled Precision =  $7 / 11 = 63.6\%$

Unlabeled Precision =  $10 / 11 = 90.9\%$

# Recall

What proportion of constituents in *gold tree* is also present in *parse tree*?



Labeled Recall =  $7 / 9 = 77.8\%$

Unlabeled Recall =  $8 / 9 = 88.9\%$

# F-Score

- Precision and recall measure opposing qualities of a parser (“soundness” and “completeness”)
- Summarize both together in the *f-score*:

$$F_1 = \frac{2 \cdot P \cdot R}{P + R}$$

- In the example, we have labeled f-score 70.0 and unlabeled f-score 89.9.

# Summary

- PCFGs extend CFGs with rule probabilities.
  - ▶ Events of random process are nonterminal expansion steps. These are all statistically independent.
  - ▶ Use Viterbi CKY parser to find most probable parse tree for a sentence in cubic time.
- Read grammars off treebanks.
  - ▶ next time: learn rule probabilities
- Evaluation of statistical parsers.